ZFS
THE LAST WORD IN FILE SYSTEMS

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ZFS Objective

End the Suffering

- Figure out why storage has gotten so complicated
- Blow away 20 years of obsolete assumptions
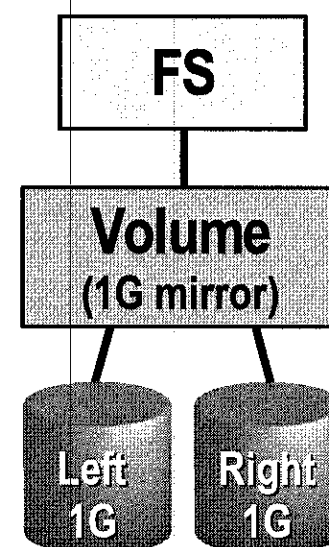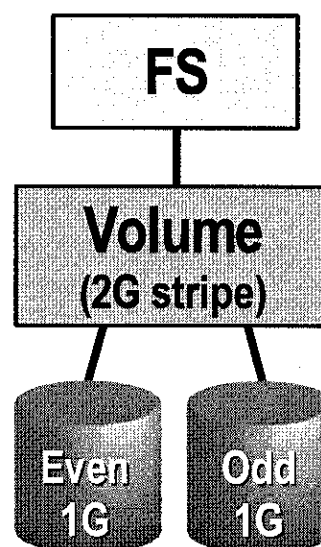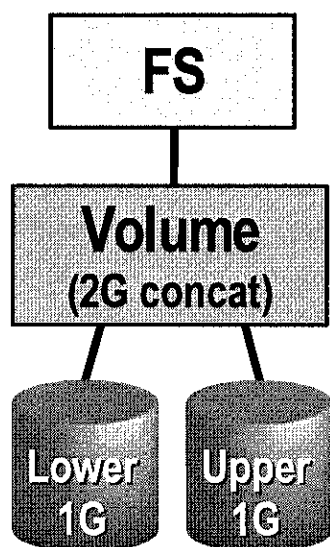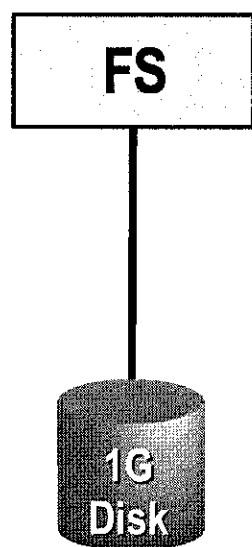- Design an integrated system from scratch

ZFS Overview

- ## Pooled storage
  - Completely eliminates the antique notion of volumes
  - Does for storage what VM did for memory

- ## Provable end-to-end data integrity
  - Detects and corrects silent data corruption
  - Historically considered "too expensive"

- ## Transactional design
  - Always consistent on disk
  - Removes most constraints on I/O order – <u>huge</u> performance wins

- ## Simple administration
  - Concisely express your intent

Sun microsystems

Why Volumes Exist

In the beginning, each filesystem managed a single disk.

- Customers wanted more space, bandwidth, reliability
  - Hard: redesign filesystems to solve these problems well
  - Easy: insert a little shim ("volume") to cobble disks together
- An industry grew up around the FS/volume model
  - Filesystems, volume managers sold as separate products
  - Inherent problems in FS/volume <u>interface</u> can't be fixed
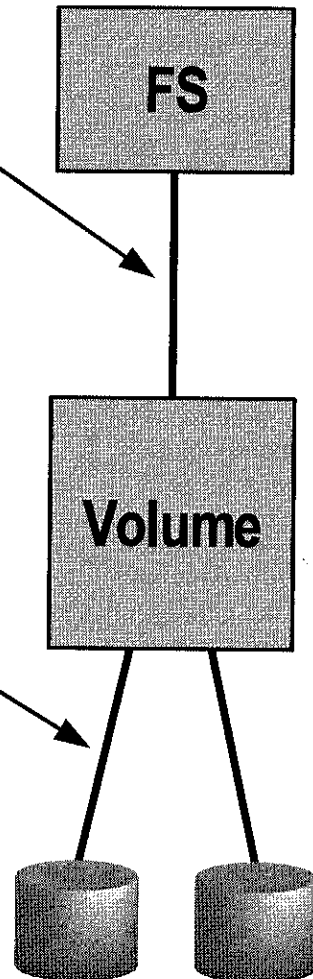
FS/Volume Model vs. ZFS

FS/Volume I/O Stack

Block Device Interface

- "Write this block, then that block, ..."

- Loss of power = loss of on-disk consistency

- Workaround: journaling, which is slow & complex

Block Device Interface

- Write each block to each disk immediately to keep mirrors in sync

- Loss of power = resync

- Synchronous and slow

FS

Volume

ZFS I/O Stack

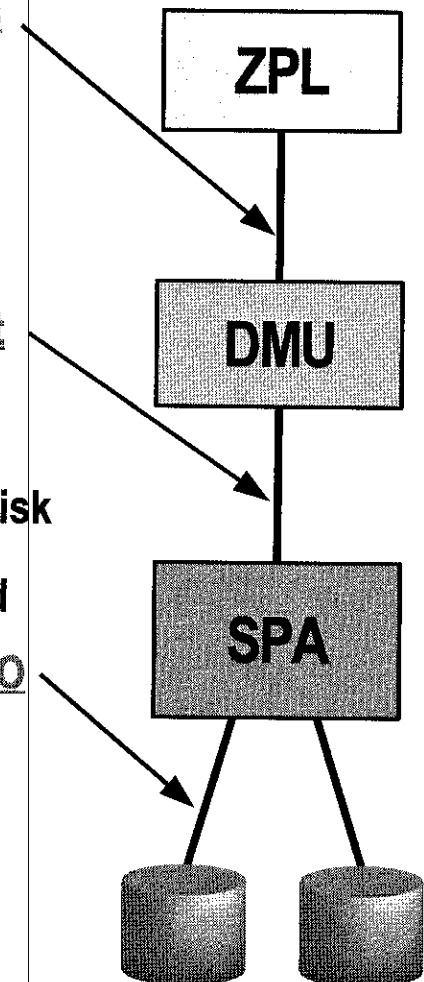**Object-Based Transactions**

- "Make these 7 changes to these 3 objects"

- All-or-nothing

Transaction Group Commit

- Again, all-or-nothing

- Always consistent on disk

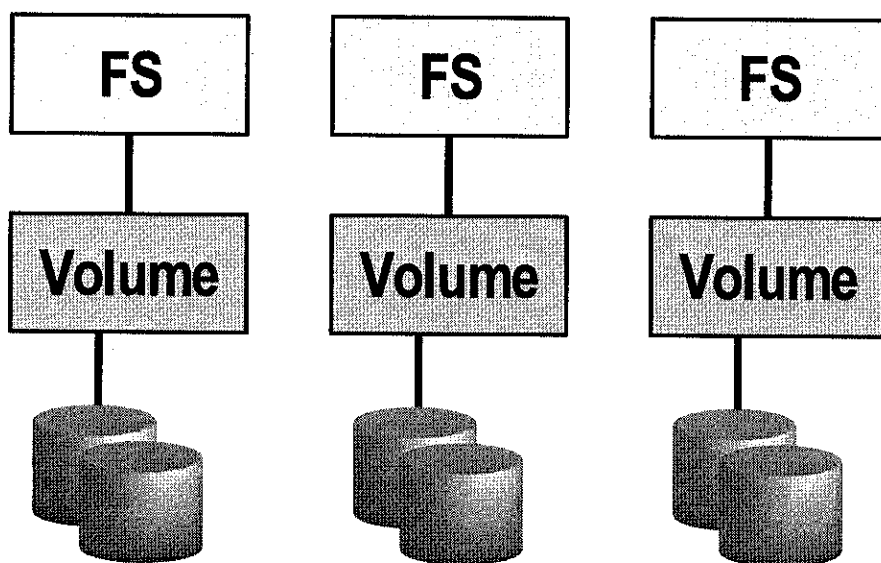- No journal – not needed

Transaction Group Batch I/O

- Schedule, aggregate, and issue I/O at will

- No resync if power lost

- Runs at platter speed

ZPL

DMU

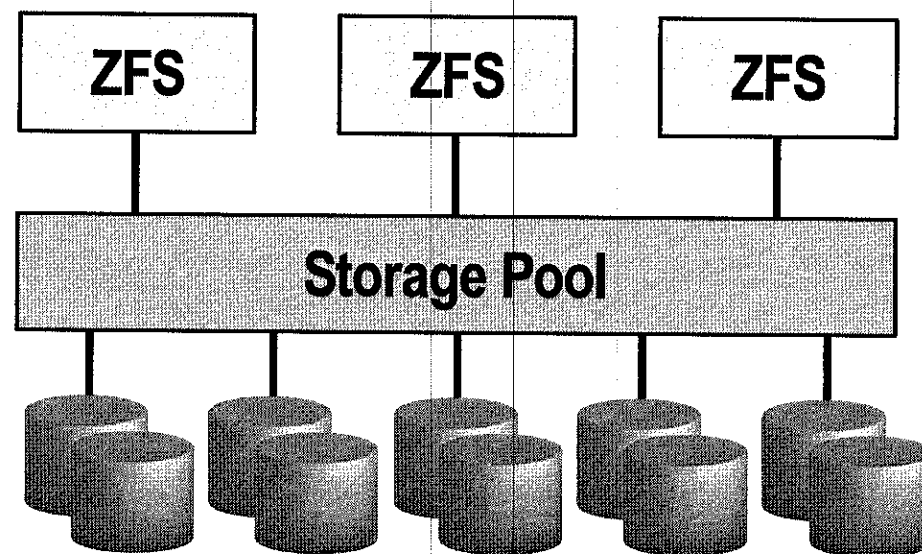**SPA**

FS/Volume Model vs. ZFS

Traditional Volumes

- Abstraction: virtual disk
- Partition/volume for each FS
- Grow/shrink by hand
- Each FS has limited bandwidth
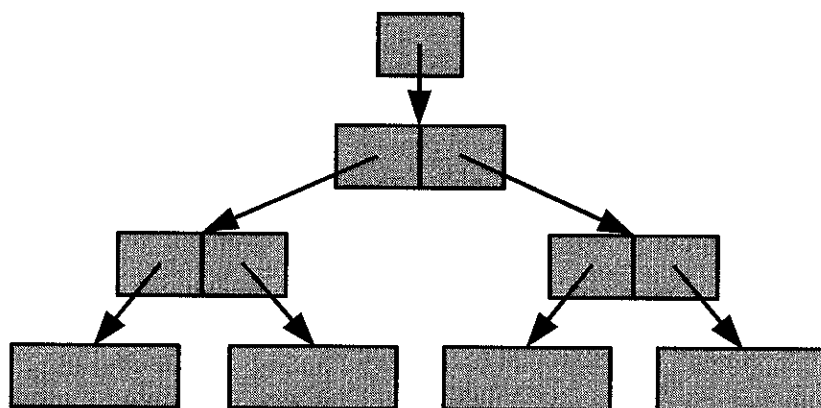- Storage is fragmented, stranded

ZFS Pooled Storage

- Abstraction: malloc/free
- No partitions to manage
- Grow/shrink automatically
- All bandwidth always available
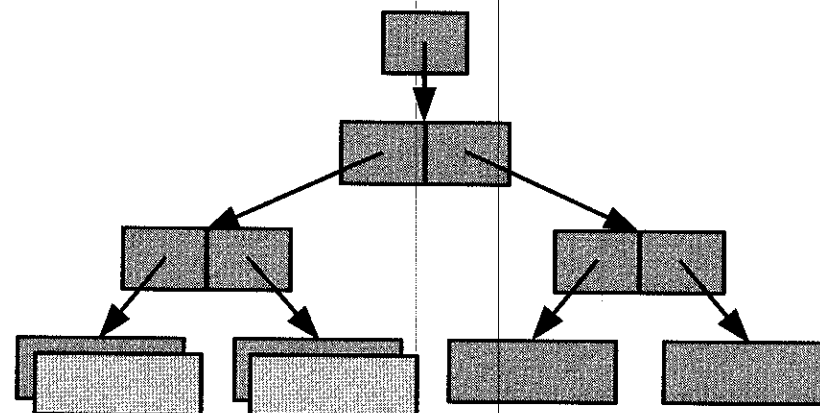- All storage in the pool is shared
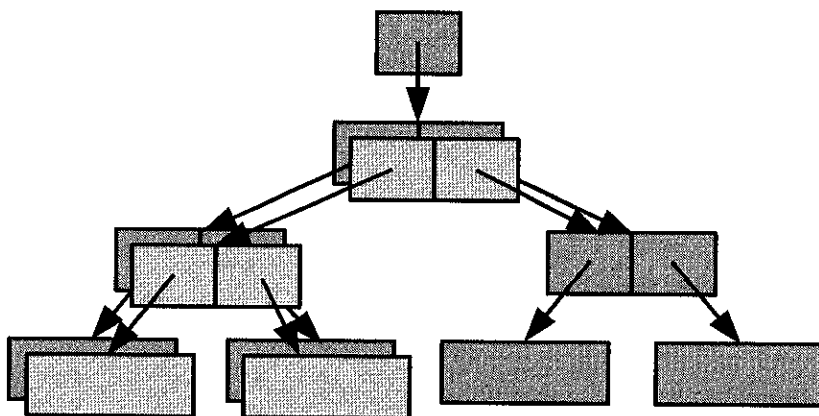
Copy-On-Write Transactions

1. Initial block tree

2. COW some blocks

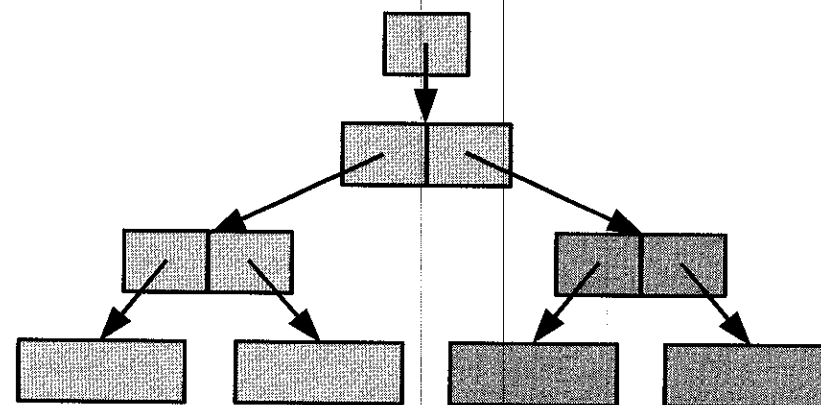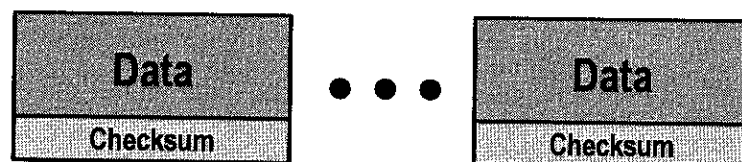

3. COW indirect blocks

4. Rewrite uberblock (atomic)

Bonus: Constant-Time Snapshots

- ## At end of TX group, don't free COWed blocks

  - ### Actually cheaper to take a snapshot than not!

Snapshot root ——————▶   ◀—————— Live root

ZFS – The Last Word in File Systems

End-to-End Data Integrity

Disk Block Checksums

- Checksum stored with data block

- Any self-consistent block will pass

- Can't even detect stray writes

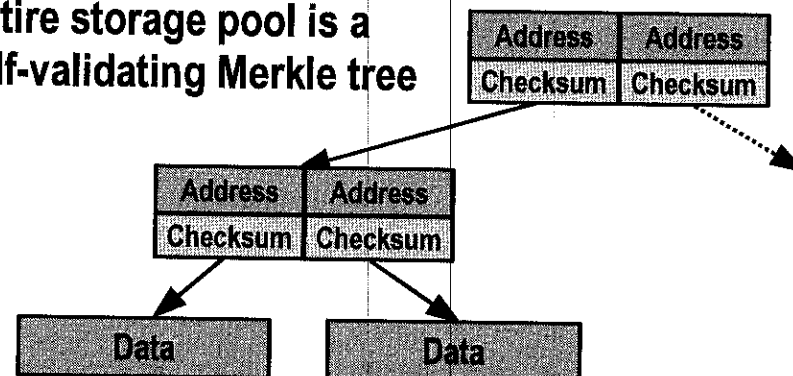- Inherent FS/volume interface limitation

ZFS Data Authentication

- Checksum stored in parent block pointer

- Fault isolation between data and checksum

- Entire storage pool is a self-validating Merkle tree

Disk checksum only validates media

✓	Bit rot
✗	Phantom writes
✗	Misdirected reads and writes
✗	DMA parity errors
✗	Driver bugs
✗	Accidental overwrite

ZFS validates the entire I/O path

✓	Bit rot
✓	Phantom writes
✓	Misdirected reads and writes
✓	DMA parity errors
✓	Driver bugs
✓	Accidental overwrite